Verification of blockchains and smart contracts Formal Methods Update, 2018 BITS Pilani Goa

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- Introduction to blockchains
- Smart contracts
- Verification issues

Introduction to blockchains

Banks and ledgers

- Record of all
 transactions
- Maintained by a
 trusted
 authority
- Each entry is
 validated
- Compute net
 balance etc

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Public ledgers

- Ledgers areprivate
- Can we maintain a public ledger?
 Eliminate trusted authority

	Jalurday
	/
Amount going of	
Total amount discounted, Bills and Notes 225.076 . 11.5	A State Section

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Challenges

 Integrity of individual
 transactions

Consensus on
 overall set of
 transactions

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A solution

- Maintain a
 distributed
 ledger
- Duplication
 prevents
 tampering
- Cryptography for
 authentication

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A solution

 A physical ledger has pages o Distributed version has blocks of data o These blocks are Linked together o Blockchain!



BLOCKS

Each block is a collection of transactions
Each block
Each block
points to parent block



Hash function

Compute random summary of input
"Impossible" to invert

o Collisions rare

Different inputs
 produce
 different outputs

0d7006cd055e94cf 614587e1d2ae0c8e

The quick brown fox jumps over the lazy dog.

Blockchain integrity

 Each block has a hash of the transactions it
 contains

Each block
 includes a hash
 of parent block



hash(parent) hash(my transactions)

Transactions

Public key cryptography

Each person P has
a public key U and
a pRivate key R
U and R are
inverses

To encrypt text t
for P to read,
send U(t)
R(U(t)) = t

The quick brown fox jumps over the lazy dog.

The quick brown fox jumps over the lazy dog.

Digital signatures

o U and R are inverses o R(U(E)) = E@ Also, U(R(E)) = E !!o Sign using R o Recipient can verify using U

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R

U

0d7006cd055e94cf 614587e1d2ae0c8e

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@ Who writes the transactions in the blockchain? No centralised authority o Transactions are created by originator

	Transaction
	From A
1000	TO B
	Amount

 A digitally signs @ Cannot repudiate later @ A uses B's public key to create a challenge only B can solve

Only B can claim
 this amount

	Transaction
	From A
	Dig Sig of A
	TO B
	Challenge
8	Amount

a Where's the money? @ No centralised authority to certify the money A holds o Must refer to previous transactions where A acquired the money

Transaction
From A
Dig Sig of A
TO B
Challenge
Amount
Sources of funds

Adding blocks

- Peer to peer
 network
- Transactions
 broadcast to all
 nodes
- Periodically,
 collect transactions
 into a block and
 add to chain



Mining blocks

- Process of adding
 a block is called
 mining
- Mining is
 decentralised
- Blockchain may
 fork
- Integrity of the ledger is lost!



Distributed consensus

All nodes should agree on blocks
Elegant solution due to Satoshi Nakomoto

Emerging
 distributed
 consensus



Proof of work

Adding a node requires solving a hashing problem
Brute force search

Calibrated so that
 it takes about 10
 minutes to solve on
 current hardware



Proof of work

After mining a
 block, miner
 broadcasts

Other miners
 abandon efforts,
 accept this block,
 move to next block

Serial numbers



Blockchain forking

o Two miners may succeed in parallel o Variants of chain may propagate Mismatch between your chain and new block - keep longer chain

@ Eventually converges

Incentive for mining

@ Why spend computational effort to mine? Transaction fees and other incentives @ Bilcoin!



Smart contracts

 A uses B's public key to create a challenge only B can solve

Only B can claim
 Chis amount

How is this done?

Transaction
From A
Dig Sig of A
TO B
Challenge
Amount

Challenge scripts

- Simple stack based programming
 Language
- Locking script
 DUP HASH160 <PubKHash> EQUALVERIFY CHECKSIG
 <PubKHash> hash of B's public key
- Unlocking script
 <Sig> <PubK>
 - Sig> <PubK> signature, public key of B

Challenge scripts...

Concatenate and execute on stack VM
 <Sig> <PubK> DUP HASH160 <PubKHash>
 EQUALVERIFY CHECKSIG

]/

More general scripts

- Multisignature
 - N public keys recorded in the script
 - M must provide signatures to unlock
- e Conditional
 - Three partners, majority must sign
 Lawyer can access with one partner

Scripting Language

o Bilcoin

- Scripting Language is intentionally
 Turing incomplete
- o Conditionals, but no loops

o Ethereum

- @ Richer Language, Turing complete
- High level language Solidity that
 compiles down to stack language

Smart contracts

- A script that executes when a
 transaction is invoked
- Ethereum contracts can express
 objects with encapsulated state
- @ Example: DAO
 - Decentralized Autonomous
 Organisation

Verification

Blockchain convergence

- Proof of work eventually convergent solution to distributed
 consensus
- @ Ensures blockchain does not fork
- Need majority collusion to fabricate
 alternate chain
 - e Would allow double spending

Vulnerability

Hijacking Bitcoin: routing attacks on cryptocurrencies, Apostolaki et al, IEEE Security and Privacy 2017

Structure of Internet is not uniform

- Concentration of switches, routers
 make partitioning possible
- Can also delay packets

Model checking

Modeling and Verification of the Bitcoin Protocol, Chaudhury et al, MARS Workshop 2015

UPPAAL model of Bitcoin network
Investigate forking, double spending
Model checking of a very small scale model, 4 nodes, 1 malicious

Smart contract verification

Online Detection of Effectively Callback Free Objects with Applications to Smart Contracts, Grossman et al, POPL 2018

- Decentralized Autonomous Organisation
- DAO bug stole \$150 million dollars
- @ Reentrant code (callbacks)

Automatic verification of effectively
 callback free objects

DAO

Object Dao
Map <Object,int> credit
int balance
Invariant
(sum o: credit[o]) = balance

Method
withdrawAll(Object o)
if (credit[o] > 0)
this.balance -=
 credit[o]
o.pay(credit[o])
credit[o] = 0

Method
deposit(Object o,
 int amount)
credit[o] += amount
balance += amount



Object Attacker

Object Dao bool stop = false int balance Method pay(int profit) this.balance +=profit if (!stop) Call stop = trueback Dao. withdrawAll(this) stop = false

Method >> withdrawAll(Object o) if (credit[o] > 0)this.balance -= credit[0] o.pay(credit[o]) credit[0] = 0Method deposit(Object o, int amount) credit[o] += amount balance += amount